

## Lecture 12 - Oct. 22

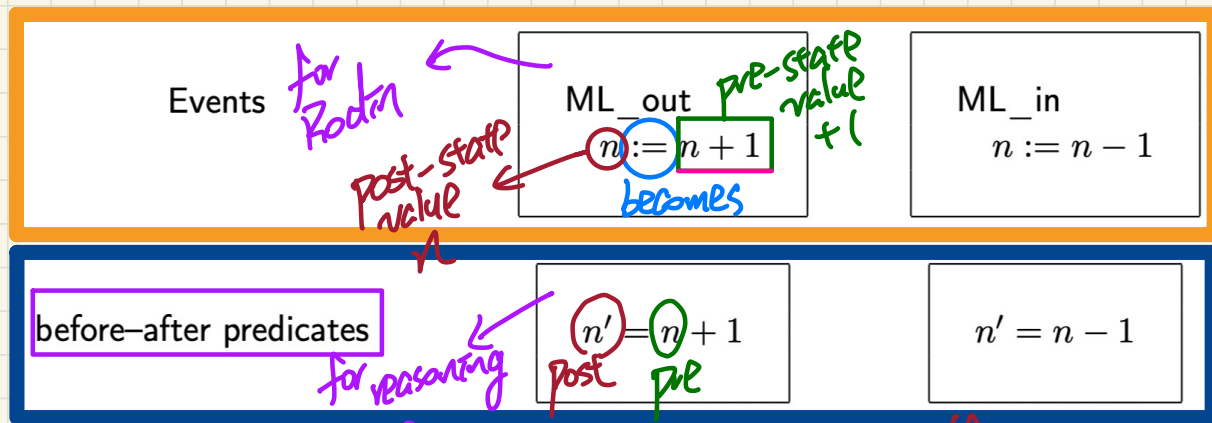
### Bridge Controller

***Event Action vs. Before-After Predicate***  
***Before- vs. After-States***  
***Sequents: Syntax and Semantics***

## Announcements/Reminders

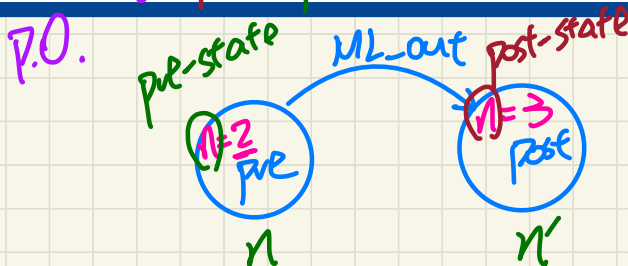
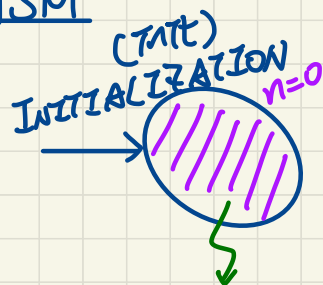
- **ProgTest1** results to be released around next Monday.
- **WrittenTest1** tomorrow during your enrolled lab session
- **Lab4** released (**ProgTest2** on November 6)
  - + Try to complete Part 1 by Friday.
  - + Follow the proof steps in Part 2 & collect questions.
  - + Scheduled lab session on **October 30**.

# Before-After Predicates of Event Actions



- Pre-State
- Post-State
- State Transition

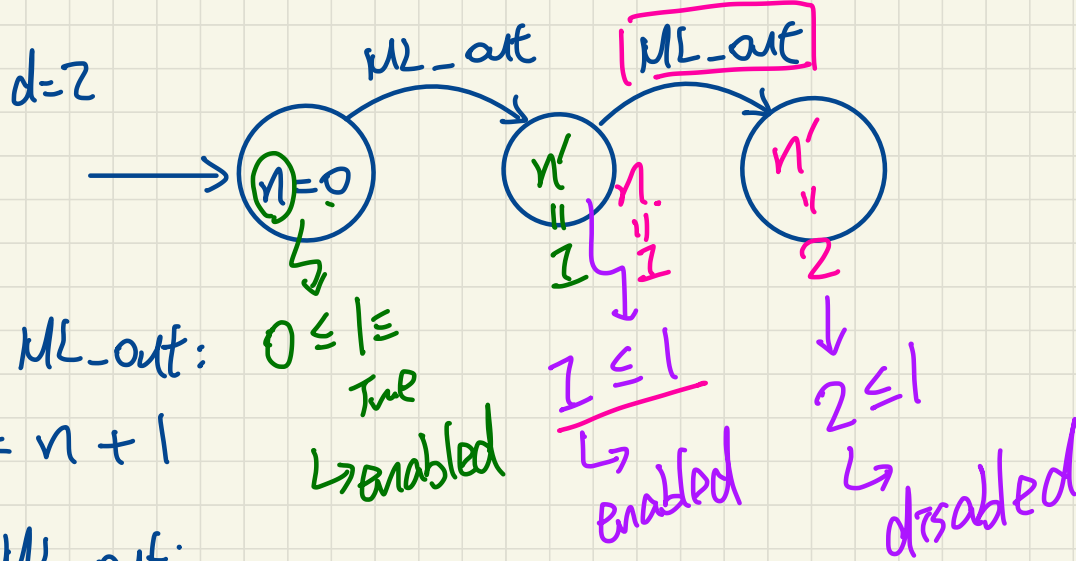
ASM



For each variable  $x$ :

- (1) Write  $x$  to denote its pre-state value.
- (2) Write  $x'$  to denote its post-state value.

$\langle \text{init} \rightarrow \text{ML\_out} \rightarrow \text{ML\_out} \rangle$



~~true say:  $n \leq 1$~~   
~~(always enabled)~~

# Transition of an Event

$$b \in \text{Account} \rightarrow \mathbb{N}$$

withdraw

$a : \text{Account}$

$v : \mathbb{N}$

where

$a \in \text{dom}(b)$

begin

$b := b \nabla \{a \mapsto b(a) - v\}$

end

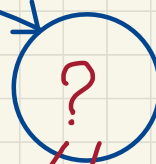
enabling condition of event  
1. guard of withdraw

2.  $I$  must be true

pre-state

withdraw

(must be enabled to occur).



post-state

Is the post-state,

after the event's action takes effect, still safe?

$\hookrightarrow$

Is  $I$  maintained?

Pre-State:

$I$

Post-State:

$I$  remains to be true

$\forall a. a \in \text{dom}(b)$

$\Rightarrow$   
 $b(a) \geq -C$

\* effect of event action.

\*

$I: \forall a. a \in \text{dom}(b) \Rightarrow b(a) \geq -C$

Invariant I in pre-state:

$$\forall x. x \in \text{dom}(b) \Rightarrow x \geq -C$$

Event  
action

$$b := b \uparrow \{a \mapsto b(a) - v\}$$

I in post-state:

$$\forall x. x \in \text{dom}(\underbrace{b \uparrow \{a \mapsto b(a) - v\}}_{\substack{\text{new value of} \\ b \text{ in post-state of withdraw.}}}) \Rightarrow x \geq -C$$

## Exercise: Event **Actions** vs. **Before-After** Predicates

Q. Are the following event **actions** suitable for a swap between  $x$  and  $y$ ?

```
swap
begin
  temp := x
  x := y
  y := temp
end
```

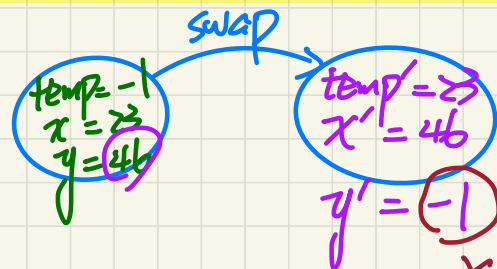
inappropriate

$:=$  should not be considered as seq. assignment

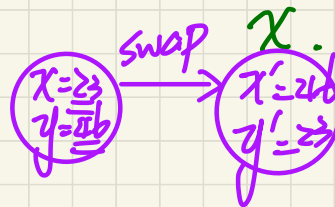
BAP:

$temp' = x$   
 $x' = y$   
 $y' = temp$

pre-state value which has no connection to



Correct:  $x := y$   
 $y := x$



BAP:  
 $x' = y$   
 $y' = x$

## Design of Events: Invariant Preservation

variables:  $n$

to be formulated as a  
proof obligation

ML\_out  
begin  
   $n := n + 1$   
end

ML\_in  
begin  
   $n := n - 1$   
end

invariants:

inv0\_1 :  $n \in \mathbb{N}$

inv0\_2 :  $n \leq d$

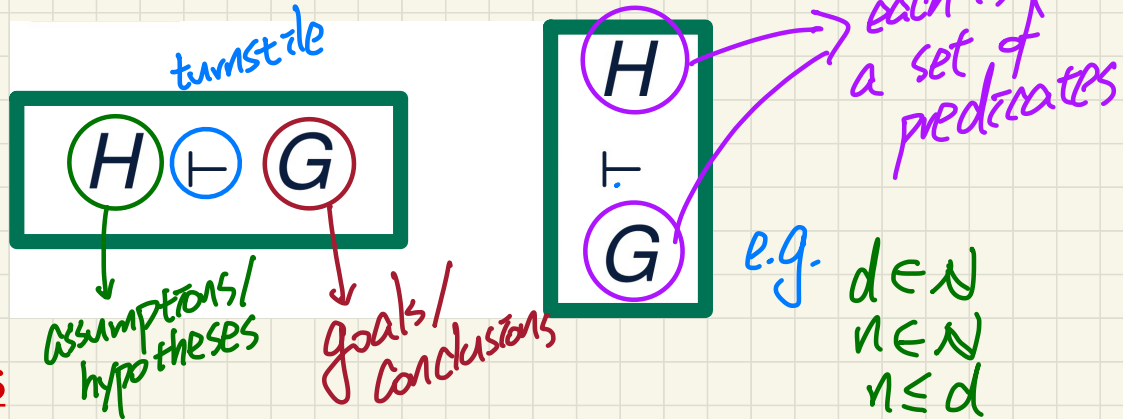
Desire:  $\forall^{n \in \mathbb{N}} \text{state} \cdot \text{state} \in \text{StateSpace}$   
 $\Rightarrow I(\text{state})$

To disprove this: find  $\text{state} \in \text{StateSpace}$  with  $n \leq d$   
but  $\neg I(\text{state})$ .



# Sequents: Syntax and Semantics

## Syntax



## Semantics

$$\frac{H \vdash G}{\text{predicate}} \Leftrightarrow H \Rightarrow G$$

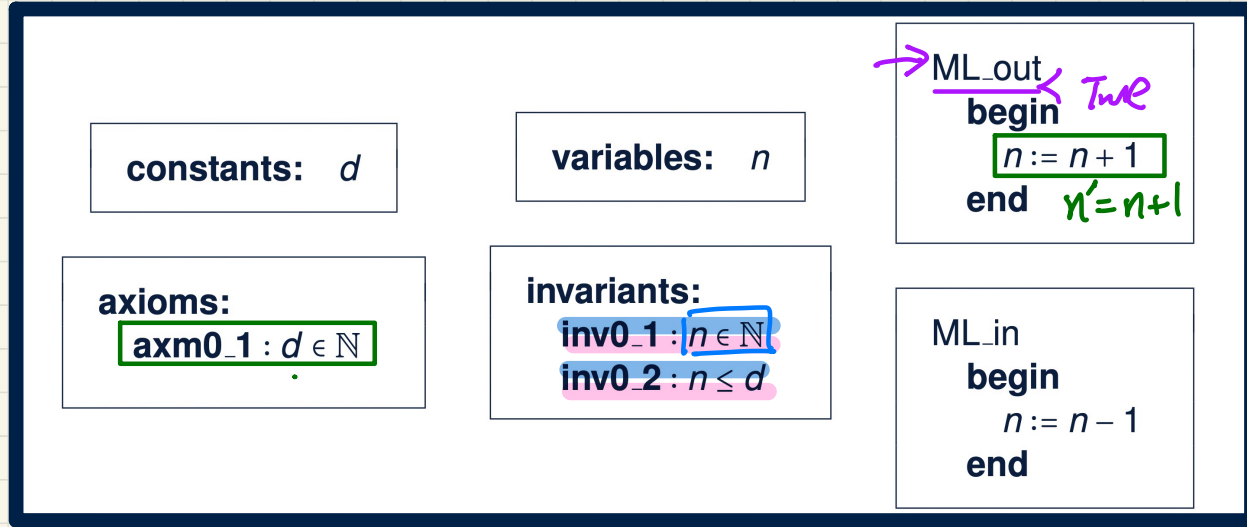
$\hookrightarrow$  assuming that  $H$  is true,  $G$  is provable.

Q. What does it mean when  $H$  is empty/absent?

- ①  $\boxed{\text{True} \vdash G} \equiv \text{True} \Rightarrow G \equiv \boxed{G}$
- ②  $\boxed{\text{False} \vdash G} \equiv \text{False} \Rightarrow G \equiv \boxed{\text{True}} \leadsto \text{nothing to prove!}$

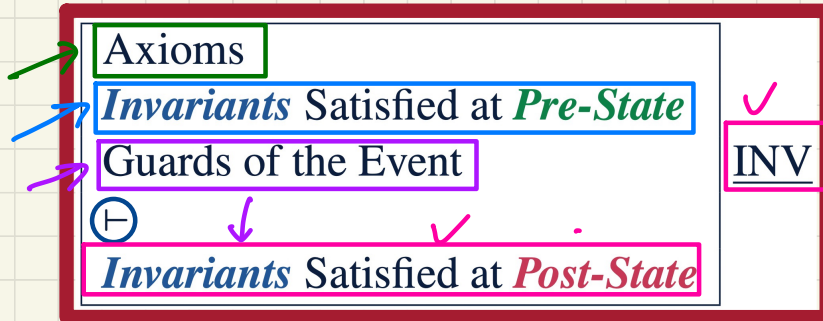
# PO/VC Rule of Invariant Preservation

Exercise ML\_in/INV



\*  $n+1 \in \mathbb{N}$   
 $\wedge$   
 $n+1 \leq d$

$d \in \mathbb{N}$   
 $n \in \mathbb{N}$   
 $n \leq d$   
 $\text{true}$   
 $\vdash$   
 $n+1 \in \mathbb{N}$   
 $n+1 \leq d$



INV

no comp of Proof obligation

$\checkmark$  ML\_out/INV

$d \in \mathbb{N}$   $\rightarrow$  pre-state  
 $n \in \mathbb{N}$   $\rightarrow$  pre-state  
 $n \leq d$   $\rightarrow$  pre-state  
 $\text{true}$   
 $\vdash$   
 $n' \in \mathbb{N} \wedge n' \leq d$   $\rightarrow$  post-state